Center for Real-Time Computing: On going efforts in parallel mesh generation and medical image computing

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CFD Vision 2030 Study: A Path to Revolutionary Computational Aero-sciences

The 90's Approach to Parallel FEM



Late 90', Cornell TC, IBM SP2 16 PEs with parallel I/O:

Mesh Size	Time (sec)	I/O	CSR Tran.	ParMetis	Data Movem.	Total (//mesh)
1M	255.9	217.8	24.28	5.46	85	588.44 (38.6) x 7
2M	461.8	471	50.58	8.58	173.3	1162.2 (85.1) x 8

Traditional			SM	GP	SMGP + ParMetis		
M_{Size}	S _{Sep.}	D _{Imb.}	S _{sep.}	D _{Imb}	S _{sep.}	D _{Imb}	
1M T	0.042	1298	0.064	3058	0.044	883	
2M T	0.038	2009	0.061	6578	0.037	1529	

I/O & data movement > 50%

Simultaneous Mesh Generation and Partitioning for 3D Domains D. Nave and N. Chrisochoides.

In Proceedings of 8th International Meshing Roundtable, Lake Tahoe, CA, pp 55-66, Sandia National Labs Publications, 1999.

Parallel FEM Simulation of Crack Propagation - Challenges, Status and Perspectives (Bruce Carter, Chuin-Shan Chen, L. Paul Chew, N. Chrisochoides, Guang R. Gao, Gerd Heber, Antony R. Ingraffea, R. Krause, Chris Myers, Demian Nave, Keshav Pingali, Paul Stodghill, Stephen S. Vavasis, P. Wawrzynek. In 7th International Workshop for Solving Irregularly Structured Problems in Parallel, Cancun, Mexico, Lecture Notes in Computer Science Vol. 1800, pp 443-449, **2000**.

The gap will only increase: I/O and Complexity & scales of models



Traditionally approach today:





30% I/O overhead W/o the cost for Data movement

and memory

contention

#Tets	Meshing	CSR	ParMetis	I/O	Total
14M	149s(1)	0.2s	7s	67s	223s

Parallel Mesh Generation today: mesh is already in place

#Tets	// Meshing	CSR	ParMetis	Ю	Total	
14M	13s(12)	0.2s	7s*	0s	20.2s	X1 ′

*Dominique Lasalle, George Karypis: Multi-threaded Graph Partitioning, IPDPS, 2013

For large scale platforms the difference is even higher!

Traditional Approach to Parallel Meshing



Parallel Mesh Generation Using an Advancing Front Method Ywith Yasushi Ito, Alan Shih, Anil Erukala, Bharat Soni, Andrey Chernikov, Nikos Chrisochoides and Kazuhiro Nakahashi. Mathematics and Computers in Simulation, Publisher Elsevier, Volume 75, No. 5, pages 200 -- 209, September, 2007



18th International Meshing Roundtable, pages 319 -- 336, Salt Lake City, Utah, October, 2009



Parallel Mesh Generation.Nikos Chrisochoides, In Numerical Solution of Partial Differential Equations on Parallel Computers,11/17/14(Eds. Are Magnus Bruaset, Petter), Springer-Verlag, 2005.

Delaunay Mesh Generation



- Bowyer-Watson kernel
 - Marked element
 - Insert its circumcenter
 - Compute its cavity
 - Triangulate its cavity

 $M_{i+1} = M_i - C_i + B_i$

- before and after BW kernel the mesh is:
 - conformal
 - Delaunay



Challenges and Requirements

Stability: the elements of the parallel generated mesh should retain the **same quality** as the elements of sequentially generated mesh;

Code re-use: meshing is labor intensive. Parallel code development still very difficult, ...

Efficiency: Single core and node performance/power

Maintenance and Correctness: Functionality, Algorithm/Software (termin.)

Code Re-use

leverage the ever evolving basic meshing algorithms/software



Parallel Optimistic Delaunay Mesh



Cavity level (i.e., medium grain for 2D, 3D and 4D: ~5, 15, 150 elems)



4D Mesh



#threads	1	32	64	128	144	160	176	1	12	
#Tets	10.7M	3.49M	7.44M	0.132B	0.151B	0.167B	0.185B	301.4K	2.837M	
Average Time(secs)	90.37	87.50	99.23	93.00	103.26	150.03	181.10	1679s	2489s	
Speedup	1.00	33.71	63.33	119.56	123.67	94.10	86.36	1	6.35	
efficiency	1.00	1.05	0.99	0.93	0.86	0.59	0.49	Overhead	871s	
								MRI Myocardium		

Real-Time Image-to-Mesh conversion for Medical Images, (with P. Foteinos): ISVD10, IMR11, MICCAI11, ACM ICS13, CGTA14, JPDC14, EwC14

Parallel Delaunay Refinement (PDR)



Sufficient condition for point Independence

Given two points p_i and p_j we can show that they can be inserted concurrently if

 $\|p_i - p_j\| \ge 4\bar{r}$

Then given a block decomposition we can concurrently refine selective blocks at a time

Practical and Efficient Point Insertion Scheduling Method for Parallel Guaranteed Quality Delaunay Refinement, (with A. Chernikov and N. Chrisochoides), In the proceedings of ACM 18th International Conference on Supercomputing, 2004.

Putting Together: PODM & PDR (cont.)

cores	1	128	144
No. elms	3.1e+7	3.74e+9	4.2e+9
T _{PODM} (s)	27.78	35.97	47.24
Speedup	1	94.5	80.8
Efficiency	1	0.74	0.56
$T_{GOSD}(\mathbf{s})$	27.78	31.46	34.04
Speedup	1	109	112
Efficiency	1	0.85	0.78



cores	1	128	144	160	176	192
No. elms	3.1e+7	3.74e+9	4.2e+9	4.67e+9	5.14e+9	5.59e+9
T _{PODM} (s)	27.78	35.97	47.24	57.24	74.20	92.77
Speedup	1	94.5	80.8	74	64	55
Efficiency	1	0.74	0.56	0.46	0.36	0.29
T _{GOSD} (s)	27.78	31.46	34.04	37.24	37.82	39.63
Speedup	1	109	112	114	124	128.2
Efficiency	1	0.85	0.78	0.71	0.7	0.67





<u>Parallel 2D Constrained Delaunay Mesh Generation</u> (with A. Chernikov and N. Chrisochoides) ACM Transactions of Mathematical Software, Vol 34(1), February, 2008.

Telescopic Approach: PCDM + LAPD (PODM + PDR)



Generalized Two-Dimensional Delaunay Mesh Refinement Andrey Chernikov and Nikos Chrisochoides. SIAM Journal on Scientific Computing, Volume 31, No. 5, pages 3387 -- 3403, 2009 Generalized Insertion Region Guides for Delaunay Mesh Refinement Andrey Chernikov and Nikos Chrisochoides. SIAM Journal on Scientific Computing, Voll. 34, No. 3, pp A1333 -- A1350, 2012

Putting It All Together

- **Code re-use:** leverage the ever evolving basic meshing algorithms/**software**
- Communication/Synchronization: leverage network/memory hierarchy

Data Decomposition

Domain Decomposition



ACM SoCG02, CGTA04, ISVD10, IMR12,CM ICS04 & ICS08, FEAD10, ACM TOMS08 [870,872] ACM ICS13, CGTA14, JPDC14a SIAM SISC11, SISC12 SIAM SISC10 JPDC14b

SIAM SISC06 SISC08

Summary

- Presented four different methods for 2D, 3D and 4D (3D+time) objects/ images and a RTS for their parallel implementation
- Map these methods according to their communication/synchronization requirements in current architectures with deep network hierarchies
- Quality is the same and sometimes better and performance is much better to known methods:

	CGAL	Tetgen	GSH3D	PODM	PODM+PDR	
Tets/sec	37K(1)	143K(1)	49K(1)	12.3M(144)	14.12(192)	A
				and guara	intee quality!	

- A plan to get to a billion-way concurrency
 - Today: 2x10² for 3D images and 10⁸ for 2D geometries
 - □ By 2017 we should have **10⁶ concurrency** (0.7 ef) for 3D images!
 - By 2020 we should be close to billion-way concurrency!