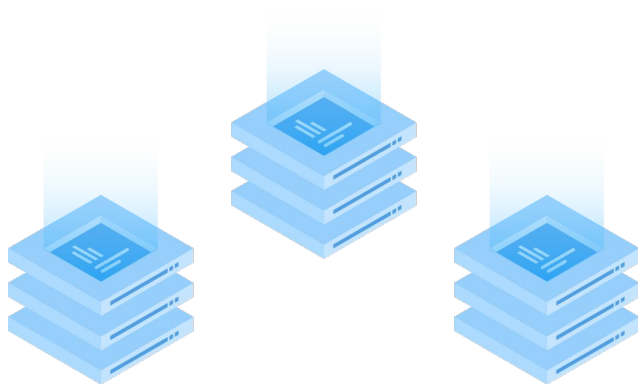


Deep Dive into the TiDB SQL Layer

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TiDB Community Slack Channel
<https://pingcap.com/tidbslack/>



About Me

- Zhang Jian (张建)
- Was: Engineer, Runtime Team, Alibaba MaxCompute
- Now: Tech Lead, TiDB SQL Engine Team
- Focus on:
 - SQL Optimization
 - Distributed Computation



Topics

1. Intro to TiDB
2. Query Optimization
3. Statistics
4. Query Execution
5. Future Work



Part I - Intro to TiDB





MySQL/MariaDB clients, ORMs, JDBC/ODBC, Applications ...

MySQL Wire Protocol

TiDB Node



TiDB Node



TiDB Node



TiDB servers, stateless, SQL engine.

Key-Value or Co-processor API calls

TiKV Node 1

Store 1



TiKV Node 2

Store 2



TiKV Node N

Store N



heartbeat

meta

PD

PD

PD

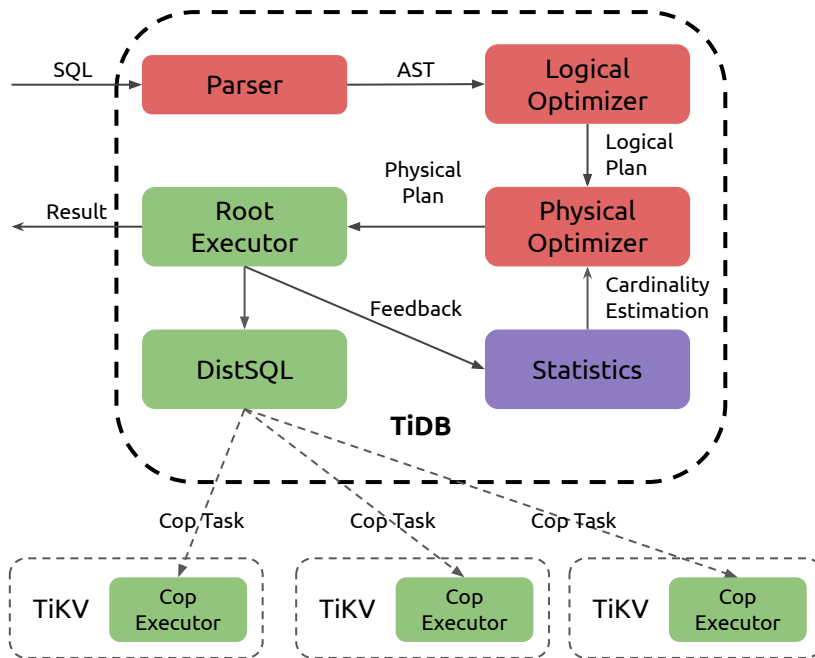
TiDB SQL Layer

Query Optimization

- Logical Optimizer
- Physical Optimizer
- Statistics

Query Execution

- Root Executor
- DistSQL
- Coprocessor Executor



Part II - Query Optimization



Query Optimization

A Query Execution Plan Example:

```
TiDB(root@127.0.0.1:test) > explain select count(*) from t1, t2 where t1.a=t2.a;
```

id	count	task	operator info
StreamAgg_13	1.00	root	funcs:count(1)
└─MergeJoin_33	12500.00	root	inner join, left key:test.t1.a, right key:test.t2.a
├─IndexReader_25	10000.00	root	index:IndexScan_24
│ └─IndexScan_24	10000.00	cop	table:t1, index:a, range:[NULL,+inf], keep order:true, stats:pseudo
└─IndexReader_28	10000.00	root	index:IndexScan_27
├─IndexScan_27	10000.00	cop	table:t2, index:a, range:[NULL,+inf], keep order:true, stats:pseudo

6 rows in set (0.00 sec)

Query Optimization

Query Optimization

- **search**, to find candidates
- **cost**, to choose the best



Query Optimization

Phase 1: Logical Optimization

- Equal
- Beneficial

```
func logicalOptimize(logic LogicalPlan) LogicalPlan {  
    for _, rule := range ruleList {  
        logic = rule.optimize(logic)  
    }  
    return logic  
}
```



Logical Rules

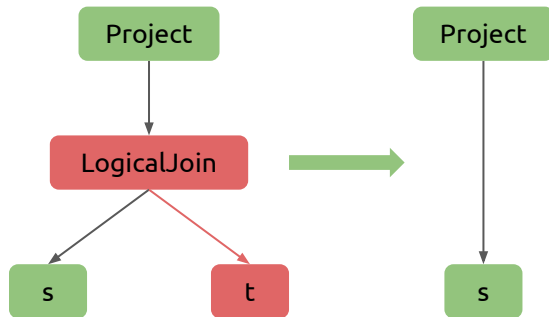
- Column Pruning
- Partition Pruning
- Group By Elimination
- Max/Min Elimination
- Project Elimination
- **Outer Join Elimination**
- Outer Join Simplification
- Subquery Decorrelation
- Predicate Push Down
- Aggregate Push Down
- TopN/Limit Push Down
- Join Reorder
- ...



Rule Example: Outer Join Elimination

- Only outer columns are used
- Join Key on inner side is unique

```
SELECT s.*  
FROM s LEFT JOIN t  
ON s.a = t.unique_key  
  
SELECT * FROM s;
```

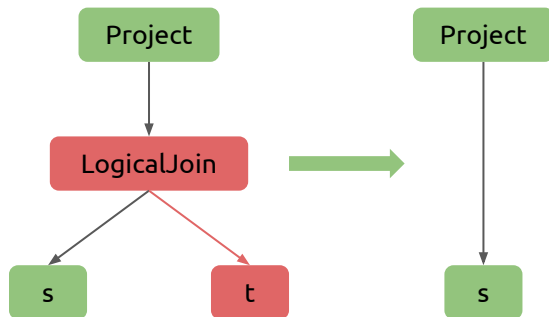


Rule Example: Outer Join Elimination

- ~~Only outer columns are used~~
- ~~Join Key on inner side is unique~~
- Only **distinct** outer columns are used

```
SELECT SUM(DISTINCT s.a)
FROM s LEFT JOIN t
ON s.b = t.b;

SELECT SUM(DISTINCT s.a) FROM s;
```



Query Optimization

Phase 1: Logical Optimization

- Equal
- Beneficial

Phase 2: Physical Optimization



Physical Optimization

Physical Property:

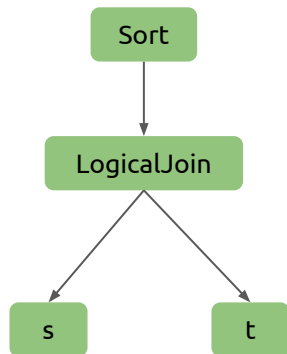
- data order
- task type

The Dynamic Programming:

- (Logical Plan , Required Physical Property) -> Physical Plan



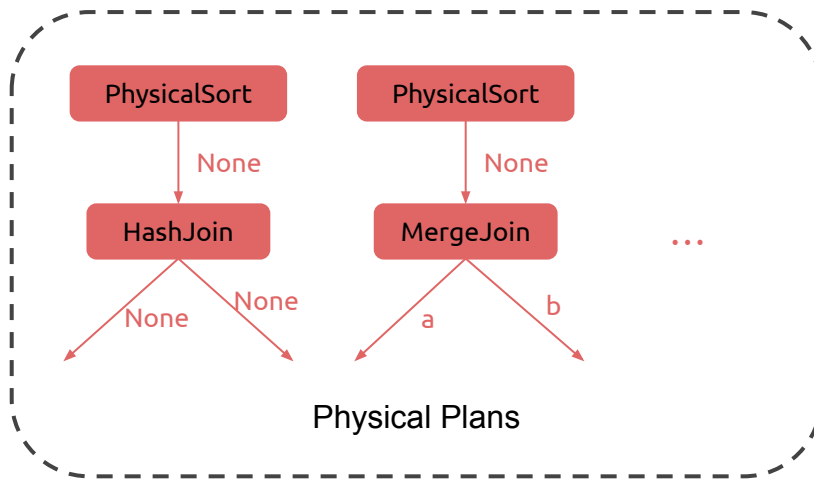
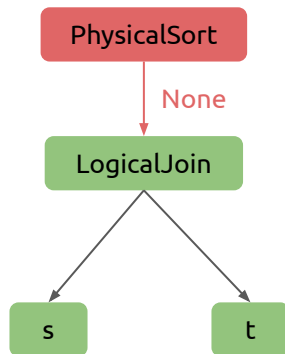
Physical Optimization



```
select *  
from s  
join t  
on s.a = t.b  
order by s.a
```



Physical Optimization



Part III - Statistics



Statistics

Cardinality Estimation

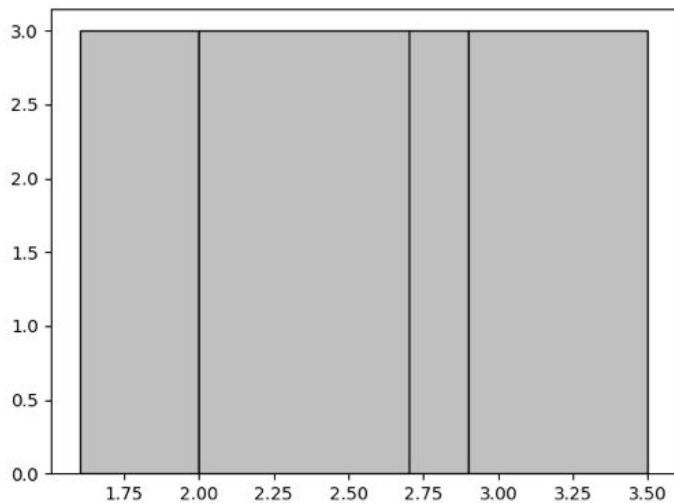
- **WHERE** $s.a > 0$
- **ON** $s.a = t.b$
- **GROUP BY** $s.a, s.b$



Statistics

What kinds of statistics does TiDB need?

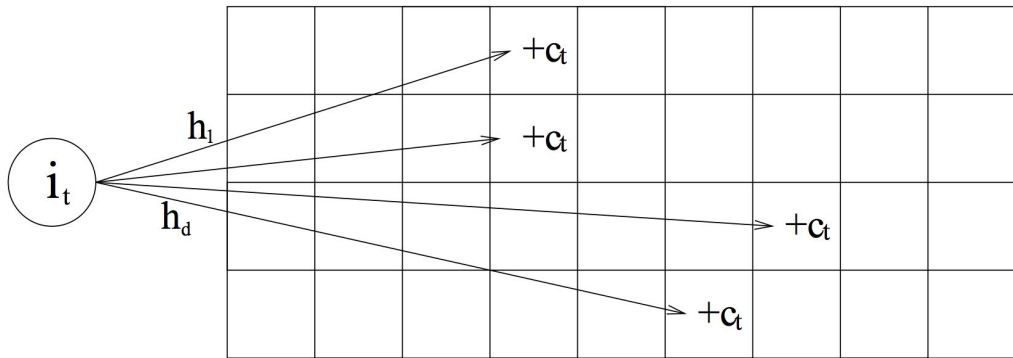
- Equi-depth Histogram



Statistics

What kinds of statistics does TiDB need?

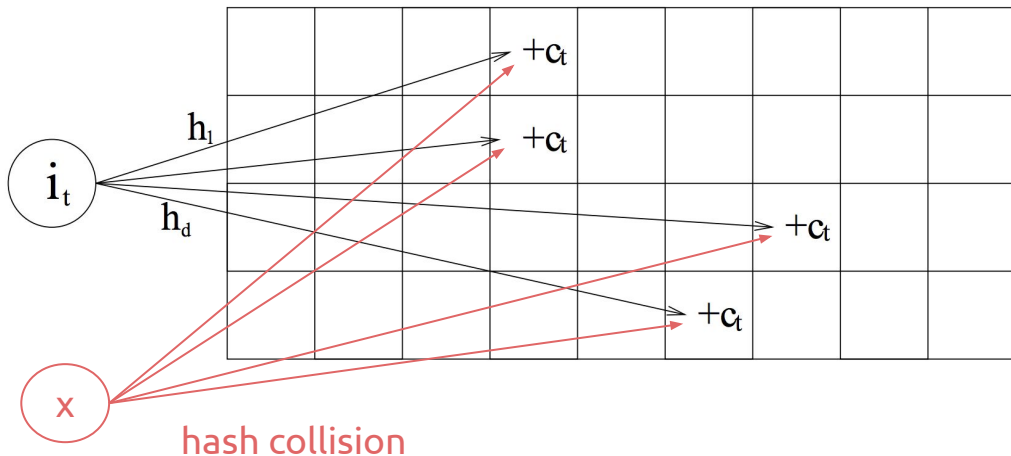
- CM-Sketch



Statistics

What kinds of statistics does TiDB need?

- CM-Sketch **with TopN**



TopN Statistic

Data	Count
x1	10000
x2	9990
...	...

Comparing to MySQL

TiDB is not a fork :-) Both systems will have unique benefits.

Major differences:

1. TiDB has more operators to consider
(hash join, merge join, hash aggregation etc.)
2. The cost model must consider network cost
3. Implements most MySQL optimizations, but currently not
index merge, loose index scan
4. Optimizer never dives the index (only considers histograms + CM-Sketch)
5. Approach is more layered (necessitated by #1)



Part IV - Query Execution



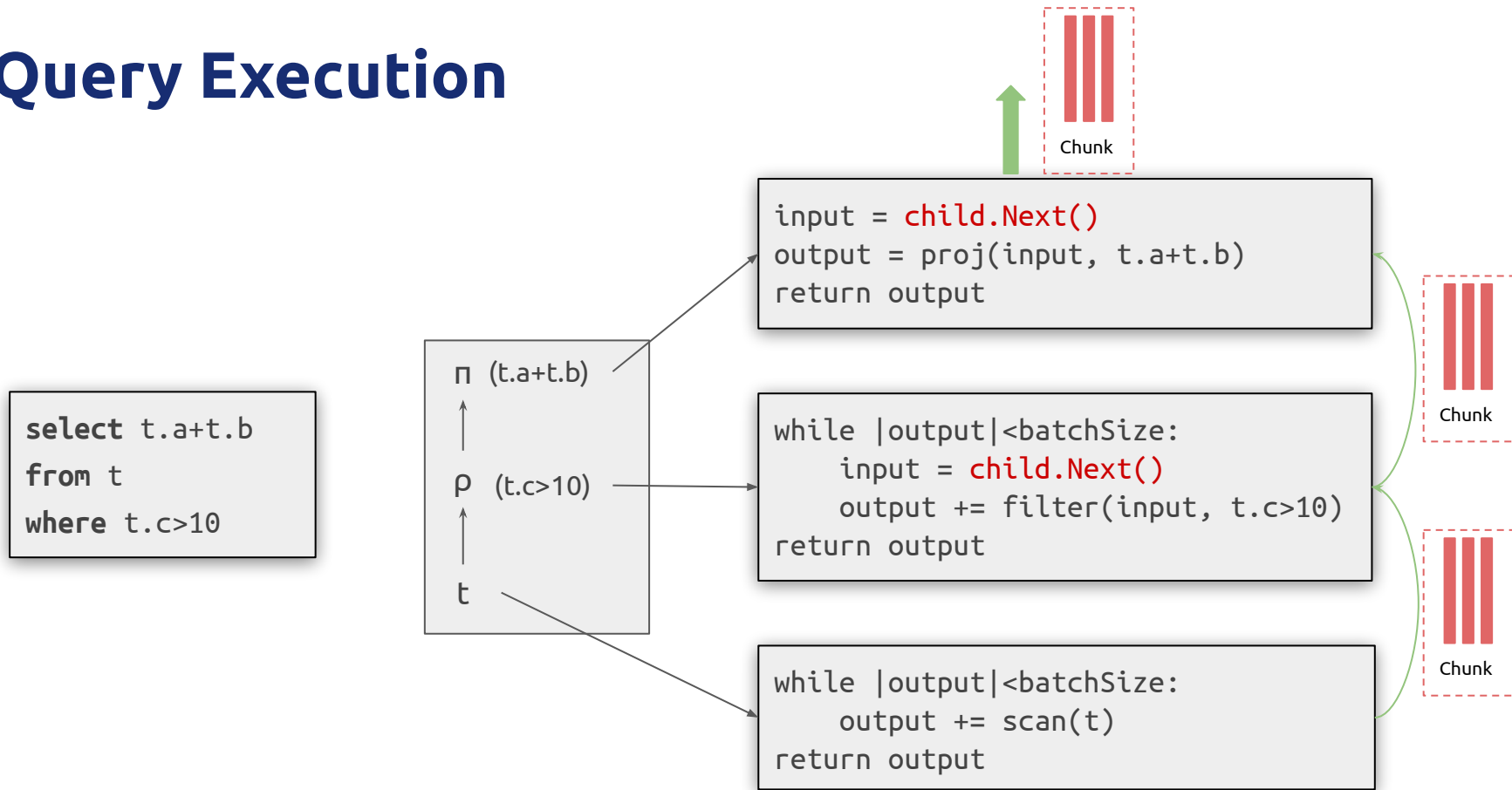
Query Execution

Since TiDB 2.0: Vectorized/Batch Model

- Based on the Iterator Model
- Column Oriented Memory Layout
- Vectorized Evaluation
- Less cache/branch/TLB miss, more efficiency



Query Execution



Query Execution

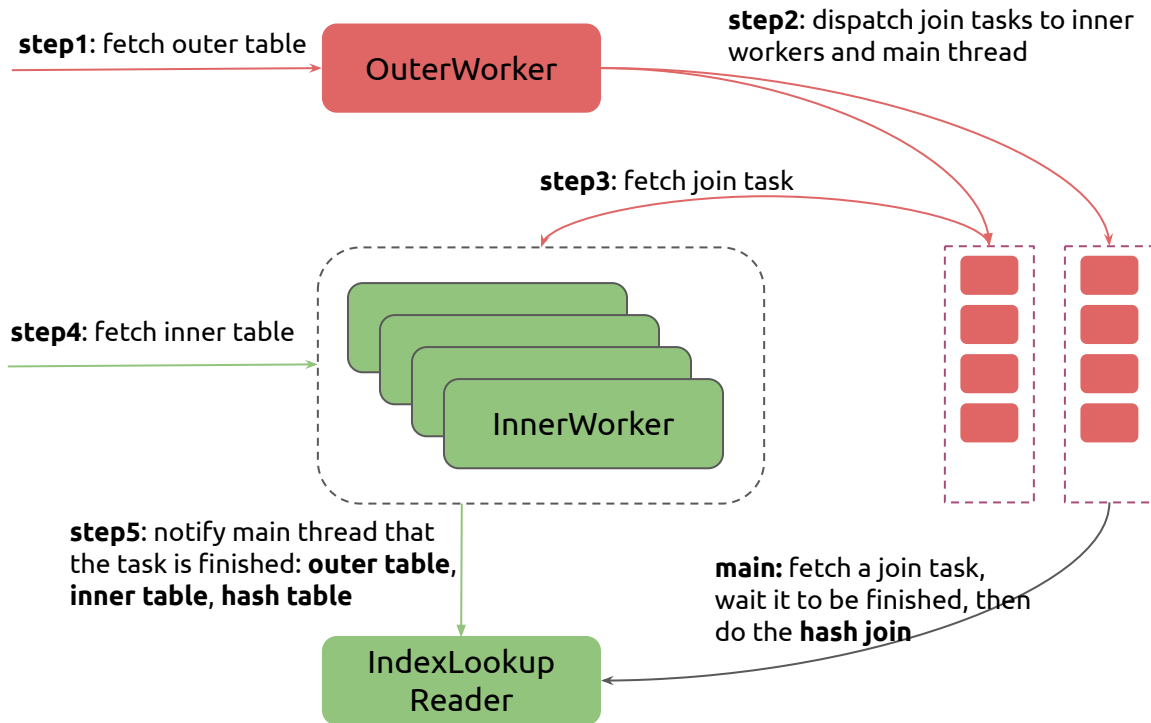
Parallel physical operator implementations:

- Hash Join
- Index Join
- Hash Aggregate
- ...

Under development: <https://github.com/pingcap/tidb/projects/11>



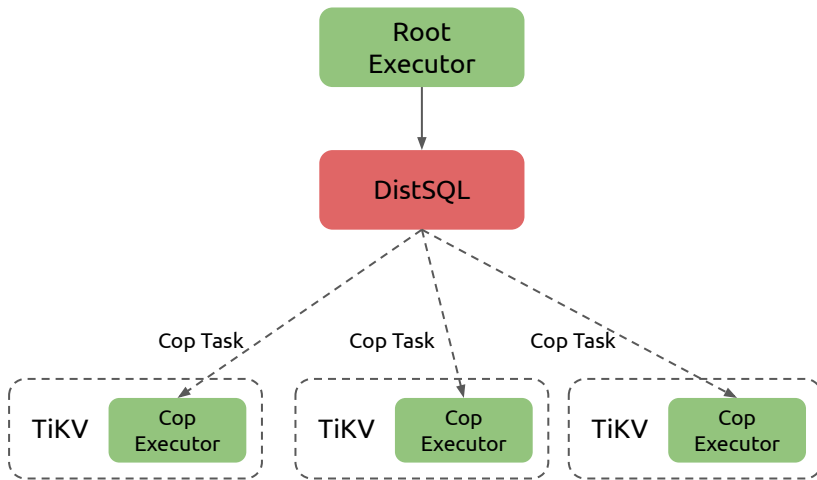
Query Execution/Index Join



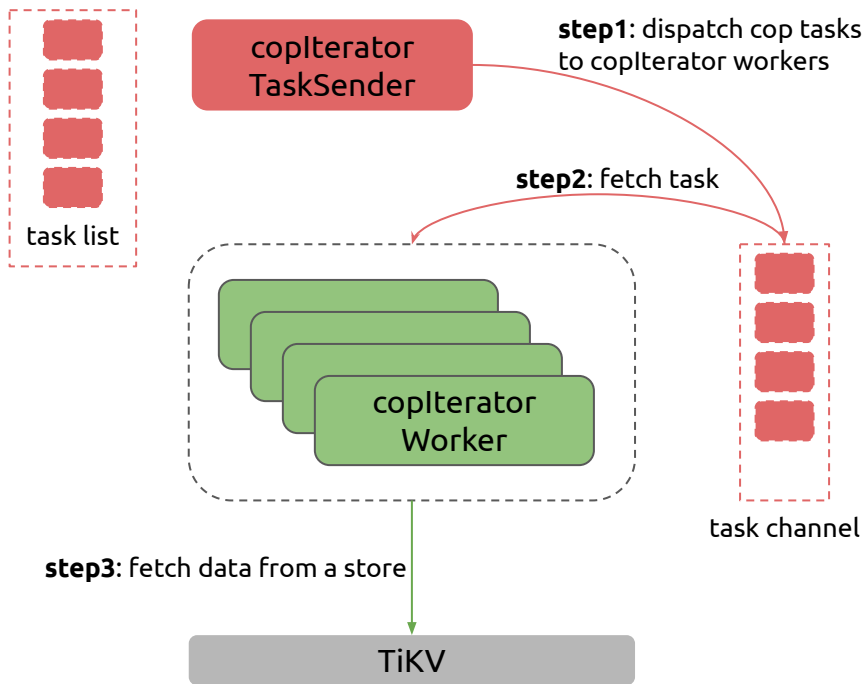
Query Execution/DistSQL

DistSQL:

1. Send Cop request
2. Receive Cop response
3. Handle TiKV Region errors



Query Execution/DistSQL



coplertator

Part V - Future Work



Future Work

Towards a Self-Driving Database

- Cascades Query Planner
- Multi-Column Statistics
- SQL Tuning Advisor
- SQL Plan Evolution
- Multi-Index Scan

Query Execution

- Parallel Control & Scheduling
- Memory Control



Thank You !



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